Analytical Results on the Performance of Shared Resource Allocation Systems

IOSIF SALEM

Division of Networks and Systems
Department of Computer Science and Engineering
CHALMERS UNIVERSITY OF TECHNOLOGY
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Department of Computer Science and Engineering
Distributed Computing and Systems Group
Division of Networks and Systems
Chalmers University of Technology
SE-412 96 GÖTEBORG, Sweden
Phone: +46 (0)31-772 10 00

Author e-mail: iosif@chalmers.se

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Iosif Salem
Division of Networks and Systems, Chalmers University of Technology

ABSTRACT

Resource allocation is a prevalent problem in a wide range of domains of computer science. Analytical tools that evaluate the performance of resource allocation systems allow us to compare with experimental ones, and utilize the design of such systems.

We consider shared-object systems that require their threads to fulfill the system jobs by first acquiring sequentially the objects needed for the jobs and then holding on to them until the job completion. Such systems are in the core of a variety of shared-resource allocation and synchronization systems. We provide methods for estimating the performance of such systems in terms of expected task throughput and delay for completion. To the best of our knowledge, this is a new perspective that can provide better analytical tools for the problem, in order to estimate performance measures similar to ones that can be acquired through experimentation on working systems and simulations.

We also study the problem of maximizing the energy utilization in the Smart Grid, where the energy supply becomes available in an online fashion (due to unpredictable energy sources) and the energy demand can have some flexibility (energy dispatch problem). Utilizing a proposed modeling of the energy dispatch problem as an online scheduling problem, we model supply-following demand in terms of the Adwords problem, in order to provide algorithmic solutions of measurable quality. In systems where demands are small compared to the individual supply, we prove a \( (1 - \frac{1}{e}) \)-competitive ratio. For cases where this does not hold, we extend the Adwords problem to utilize dynamic budgets, and present an algorithm with a \( \frac{1}{2} \)-competitive ratio.

Keywords: Resource allocation, analytical performance evaluation, shared-object systems, energy dispatch problem.
Preface

Parts of the contributions presented in this thesis have previously appeared in the following manuscripts.


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Part I

INTRODUCTION
Introduction

Resource allocation is an often challenging problem appearing in a diverse range of areas in computer science, but also in other domains such as economics and project management. In computing systems, any running process needs to use resources such as shared memory, I/O, CPU time etc., and their allocation affects not only the process’ execution but also the execution of other processes and the whole system’s performance [1, 2]. Focusing on the process interaction and to improve resource utilization, we need these shared resources to be used (alt. accessed) concurrently by many processes. Nevertheless, each process’ access to resources should not compromise the result of other processes’ executions and the communication among processes for sharing these resources should be done in a correct and efficient manner [3].

The current and future electric power grid is another example of resource
allocation systems. In such a system, that is often referred to as the Smart Grid [4], the energy supply originates from the utility company, but also from renewable energy sources (e.g. photovoltaic arrays, wind generator farms). Furthermore, energy loads can be distributed or may have less regular patterns compared to households (e.g. electric car fleets, data centers), and on the other hand, the demand for energy can have some flexibility. In this context, scheduling the allocation of energy supply to match the demand determines not only the performance on a system level but also the utilization of the system’s available resources [5]. Focusing on these domains, this thesis includes two papers (Paper I and Paper II of Part II) that contribute in the area of analytically evaluating the performance of shared resource allocation systems.

The remainder of this section is organized as follows: in Sections 1.1 and 1.2 we explain the area and the challenges for Paper I, and respectively, Paper II. In Section 1.3 we highlight the contributions of these papers and in Section 1.4 we conclude and discuss future work.

1.1 Shared Resource Allocation: Models and Performance Analysis

Modeling a shared resource allocation system (SRAS) allows us to draw conclusions on its performance but it can also utilize the design of such systems. The performance of a SRAS is commonly evaluated either through experimentation, or simulation, or analytically. We focus on a class of SRAS in which computing entities that we refer to as threads carry out tasks (jobs) that require mutual exclusion to a number of shared resources (objects). These jobs arrive at the system in a given rate and are then assigned to the threads. Our work in Paper I of this thesis contributes to the analytical performance evaluation of such systems.

We study an SRAS model that includes $N$ threads, $M$ objects and $J$ jobs. A job is a tuple including a vector of objects and an operation to be executed on those objects. Jobs are continuously arriving in the system and are then as-
1.1. MODELS AND PERFORMANCE ANALYSIS

signed to threads in a FIFO order. A thread carries out a job by (1) acquiring and holding on to each object in the job’s object vector and (2) executing the job operation after having acquired all of the job’s objects. Note that a thread acquires the job’s objects in the order of the job’s object vector and we require the objects in these vectors to follow an ascending order for the absence of deadlocks, i.e., to ensure that the delay of a thread to complete a job is bounded. Moreover, the competing threads acquire each object in a FIFO manner. Therefore, the delay for a thread to complete a job in such a system depends on that of other threads and so on. Furthermore, the contention caused due to these dependencies varies with the rate in which jobs arrive in the system (job arrival rate). We say that such a system is in equilibrium, when the rate of arriving jobs matches the rate in which these jobs are completed (completion rate). Paper I contributes to the identification of these dependencies and gives tools for analytically evaluating the performance of such systems in equilibrium.

The existing practice considers job delay and completion rate as the performance measures of working systems. Empirical experiments often study shared resource systems at their saturation point in which the system is at its peak utilization. Let us describe peak utilization scenarios using two vectors; one for job arrival rates and another for their completion rates. A saturation point is the case in which: (1) the system is in equilibrium, i.e., the arrival rate of any particular job matches the completion rate of this job, as well as (2) the system is at the stage at which a higher arrival rate of any job to the system cannot increase the completion rates. Our study considers the entire range of these equilibria rather than just peak utilization scenarios. We then propose a procedure for finding such equilibria in an approximated fashion, if such exist in the given system. Once we approximate an equilibrium, we can estimate its performance measures, i.e., job delay, completion rate and blocking time.

This problem has well-known results studying the worst-case job delays, which may even be exponential on metrics, such as the chromatic number of the resource graph [6, 7]. In this graph, the vertices (objects) are connected if there is at least one thread that may request them both at any point in time. In the context of working systems, the expected time is rather different than the worst
case and therefore computer experiments are the common way for evaluating the system performance. A generalization of the dining philosophers problem, as in [7, 8], in which every job includes a fixed set of objects that it may need is also related to this problem. We provide a new perspective that enables an analysis of the evaluation metrics by considering measures both at the system level and at the level of each resource. In particular, we consider performance measures that are associated with each resource, such as the delay, completion rate and blocking time. On the system level, we consider the job arrival and completion rates, as well as the total number of threads, $N$, and objects, $M$.

1.2 Shared Resource Allocation: The Smart Grid Case

The supply and demand in the electric power grid is a real-life paradigm of a resource allocation system. Until recently the utility company was the only supplier of energy to the electric power grid and energy should be always available on demand. Nowadays, the energy mix includes also renewable energy sources (e.g., photovoltaic arrays, wind generator farms) and possibly storage, i.e., battery arrays. Moreover, the needs from the consumers’ side are also changing, since energy loads may have less regular patterns compared to households (e.g., electric car fleets, data centers). Thus, the supply and demand in the electric power grid (often referred to as Smart Grid [4]) has began shifting towards a market-oriented paradigm, where generated energy comes from many different sources and is brokered to consumers through utilities or electricity vendors.

In Paper II of this thesis, we address the resource allocation problem of utilizing the available energy supply through flexibility in demand (or supply-following demand [9]). We provide two online algorithms that can solve this resource utilization problem in a variety of supply-related assumptions (e.g. availability of storage) along with an analysis of their competitive ratios [10]. In this direction, we connect our analysis of the energy utilization problem with the Adwords problem [11]. Paper II also provides an experimental study that
evaluates the algorithmic solution against real consumption data from a pilot housing project.

A number of approaches in the literature exist regarding the problem of supply-following demand. In [12], Kok et al. focus on both the supply and demand, and use a hierarchical mechanism and a market structure to match consumers with producers, with the ultimate goal of reducing peaks in consumption. On the other hand, Barker et al. [13] focus on background consumption loads (i.e. loads that the consumer does not interact with), and by applying scheduling techniques, such as a variation of the Earliest Deadline First algorithm, they shift demand during the day in order to reduce peaks in consumption. The works of Lu et al. [14] and Tu et al. [15] are closer to the context of Paper II of this thesis, since both present online algorithms with proven competitive ratios. However, they do so for special cases of interest: in [14], Lu et al. focus on fast-responding generators (e.g. gas or diesel turbines) and present an algorithm that operates for any combination of demand, supply and price, and in [15], Tu et al. focus on data centers and on a cost minimization problem where price is a parameter. In addition, Georgiadis et al. [16] present a novel modeling and an online algorithm that can schedule flexible demand in order to reduce peaks in consumption in scenarios where forecasts are unreliable or not available (e.g. renewable energy sources, energy storage). Nevertheless, one common element of all approaches above is that the optimization goal is the reduction of peak demand in the considered time period, while the criterion in this work is utilization of all available supply\(^1\).

## 1.3 Contributions

The contributions of this thesis lie in the area of providing analytical results for the performance of shared resource allocation systems. In Paper I of this thesis, we provide methods for analytically estimating the expected task throughput and delay for completion in a model of shared resource allocation systems\(^1\).

\(^1\)I.e. a peak might be desirable in a specific time where supply is too high, e.g. due to increased wind generation.
where computing entities require mutual exclusion to a number of shared resources due to continuously arriving tasks that they carry out. We also provide a rigorous analysis of the proposed algorithm using tools from Queueing Theory and especially Queueing Networks [17]. In Paper II of this thesis, we provide two algorithms for approximating the optimal utilization of the energy supply with respect to the flexibility of demand for the energy mix in the Smart Grid [4]. Paper II also includes two online algorithms along with proofs of their competitive ratios.

1.4 Conclusions and Future Work

This thesis provides analytical tools for evaluating the performance of resource allocation systems in two directions. In Paper I we study a system model where computing entities require mutually exclusive access to a number of resources, due to continuously arriving tasks that they carry out. The results of Paper I can be used as a pillar for modeling the performance of more complex systems, where different mechanisms are used for allocating resources, rather than the sequential acquisition that is studied here. In Paper II we provide two online algorithms for the problem maximizing the energy utilization in the Smart Grid. A possible continuation of this work could study the interaction of various pricing mechanisms with the proposed algorithms and the combined impact on maximizing resource utilization. The discussed papers follow in Part II of this thesis.

Bibliography


