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ABSTRACT

Multiple Facets (MF) is a dynamic enforcement mechanism which has proved to be a good fit for implementing information flow security for JavaScript. It relies on multi executing the program, once per each security level or view, to achieve soundness. By looking inside programs, MF encodes the views to reduce the number of needed multi-executions.

In this work, we extend Multiple Facets in three directions. First, we propose a new version of MF for arbitrary lattices, called Generalised Multiple Facets, or GMF. GMF strictly generalizes MF, which was originally proposed for a specific lattice of principals. Second, we propose a new optimization on top of GMF that further reduces the number of executions. Third, we strengthen the security guarantees provided by Multiple Facets by proposing a termination sensitive version that eliminates covert channels due to termination.

ACM Reference Format:

1 INTRODUCTION

JavaScript has become the de facto programming language of the Web. Web browsers daily execute thousands of JavaScript lines which usually have access to confidential information, for example cookies that mark that the user in a web session is authenticated. It is not surprising that JavaScript is a common target for attacks. While browsers deploy security measures in the form of access control (e.g., SOP and CSP), they are insufficient [12, 17, 30] to protect confidentiality of data.

Information flow control (IFC) is a promising technology which provides a systematic solution to handle unintentional or malicious leaks of confidential information. Recently, dynamic IFC analyses have received a lot of attention [1–3, 5, 7, 9, 10, 14, 26, 33], due, in part, to its applicability to JavaScript—where static analyses are rather an awkward fit [29].

In order to scale, a suitable IFC technique for the web not only needs to be dynamic but also needs to reduce to the minimum the modifications required to existing JavaScript code. In this light, an interesting dynamic IFC technique which fulfills both of these requirements consists in executing several copies of a program: one execution per each security level or view. In that manner, each copy of the program (view) depends only on information observable to the corresponding security level, where no leaks are therefore possible. Secure Multi Execution (SME) [14] and Multiple Facets (MF) [3] are two techniques based on this idea.

Both techniques have been proved to be a good fit for information flow security in the web since they have been successfully implemented as extensions of the Firefox browser [13, 33]. Although both SME and MF are based on multi-executions, they present important differences [7]. On one hand, SME is black-box [24], i.e., it is a mechanism that does not look inside programs but rather change the semantics of inputs and outputs to ensure security. For a moment, we assume a scenario where security levels are simply sets of principals (e.g., web origins) which denote those authorities with confidentiality concerns over data. In such a scenario, SME needs to spawn one execution for any possible set of principals—where the number of executions grows exponentially with respect to the number of principals! Instead, MF [3] is designed to reduce the number of multi-executions and the memory footprint of SME. It does so by inspecting programs code and multi-executing instructions and multiplexing memory only when needed. While MF is more resource-friendly than SME, SME provides stronger security guarantees when it comes to leaks via abnormal termination [7].

Our broad goal is to augment the efficiency of techniques based on MF and SME to general cases. In particular, we discovered that MF might sometimes spawn more multi-executions than SME—something that is counter-intuitive when considering the purpose of MF (see Section 2). Our first contribution consists on a novel technique to further reduce the number of multi-executions (and memory footprint) of MF. Our second contribution is to generalize MF to work for arbitrary lattices (see Section 3) rather than being restricted to the security lattice of principals as in the original proposal [3]. This becomes useful when, for instance, a program depends on 5 security levels. In such case, as stated originally, MF will need to encode them by using (at least) 3 principals (23 > 5), and thus execute the program 8 times, while SME will execute it only 5 times (one per security level). Finally, we combine MF and SME into a single new dynamic IFC mechanism in order to provide security guarantees as strong as SME (i.e., termination sensitive...
non-interference) while avoiding multi-executions as much as our optimized version of MF allows it. All proofs can be found in [23].

2 BACKGROUND ON SME AND MF

In this section, we discuss how on one hand, the underpinning mechanism in MF reduces the number of executions compared to SME, and on the other hand, may run more multi executions than SME because of the security lattice based on principals. Our goal here is partly pedagogical and partly to motivate and provide intuition on the optimization proposed in Section 4.

Language and Semantics To investigate the foundation of multiple facets, we use a simple, deterministic while language. Its syntax includes programs P, variables x, expressions e, and values v. We use the symbol ⊕ for binary expression operators. A value is either an integer value or a boolean value.

Figure 1 presents standard big-step semantics of the language. Memories µ map variables to values; we overload the notation of memory and use µ(e) as the evaluation function for expression e in memory µ, where µ(ν = v) = v and µ(e1 ⊕ e2) = µ(e1) ⊕ µ(e2). We write (P, µ) ⊑ P to mean that the evaluation of program P on memory µ terminates with memory µ′. We use µ[x → v] for the memory µ′ where µ′(y) = µ(y) if y ≠ x, and µ′(y) = v if y = x.

MF may use fewer resources than SME SME [14] multi executes programs, in a blackbox manner, as many times as security levels in a lattice. Let’s define an SME memory as a function that maps each variable to an array of values, one value per security level. For the sake of simplicity, let’s consider first a security lattice with only two elements H and L where H ⊈ L is the only disallowed flow. Thus, an SME memory µ maps variables to an array of 2 (possibly different) values: one corresponding to the H view and one corresponding to the L view. Let’s denote such array of values as µ(x) = (µ1(x), µ2(x)), where µ1(x) is a private, H, view and µ2(x) is a public, L, view. Assume that H(µ) (resp. L(µ)) is a memory in the standard semantics, obtained by projection of µ, mapping variables to single values of the high view (resp. low view). Then, the SME monitoring rule1 for such a language can be given by the relation µ ⊑ SME-TINI as follows:

\[
\begin{align*}
\text{SME-TINI} & \quad \begin{array}{l}
(P, H(µ)) \downarrow \mu_1 \quad (P, L(µ)) \downarrow \mu_2 \\
(P, \tilde{µ}) \downarrow \Sigma_{\text{SME-TINI}} \mu_1 \odot \mu_2
\end{array} \\
\text{SME-optim} & \quad \begin{array}{l}
(P, H(µ)) \downarrow \mu_1 \quad (P, L(µ)) \downarrow \mu_2 \\
(P, \tilde{µ}) \downarrow \Sigma_{\text{SME-TINI}} \mu_1 \odot \mu_2 \\
\end{array}
\end{align*}
\]

where ⊕ combines two normal memories into a SME memory in such a way that H(µ1 ⊓ µ2) = µ1 and L(µ1 ⊓ µ2) = µ2. The SME mechanism will blindly execute the program as many times as possible views (or positions of the array) may exist.

Consider a program h := l where initial views for variables l and h are given by: µ(h) = (1 : 0) and µ(l) = (1 : 1). In SME, using the SME-TINI rule, the assignment will be executed twice: once with H(µ) = [h → 1, l → 1] for the high view and once with L(µ) = [h → 0, l → 1] for the low view. After execution, the final SME memory will map h to (1 : 1). One way to reduce the number of executions is to exploit the knowledge that the high and the low view for variable l are equal, i.e., H(µ) = L(µ). Since the semantics is deterministic, there is no need to execute the program twice. We can use this knowledge by specialising SME at the granularity of commands and include the following assignment rule:

\[
\begin{align*}
\text{SME-optim} & \quad \begin{array}{l}
(P, H(µ)) \downarrow \mu_1 \quad (P, L(µ)) \downarrow \mu_2 \\
(P, \tilde{µ}) \downarrow \Sigma_{\text{SME-TINI}} \mu_1 \odot \mu_2 \\
\end{array}
\end{align*}
\]

Notice that this SME optimization requires to look inside the shape of the program to evaluate if expression e of an assignment satisfies the hypothesis.

In general, in order to reduce the number of executions using the multi-execution technique of SME-TINI, it is sufficient to (i) identify in an SME memory which values in the array of values are equal and (ii) remember which values correspond to which views. MF uses the multi-execution technique, implements (i) and (ii) and hence, reduces the number of executions. MF encodes values in SME memories (arrays with as many positions as lattice elements) as ordered binary trees, where the order is given by the elements of the lattice. For example, for a SME memory where µ(h) = (1 : 0 : 0 : 0) for a lattice of 4 elements with top element T, an equivalent MF memory encodes this array as (⊤1 : 0 : 0 : 0) with the meaning that 1 is the view for T and 0 for the rest. Every execution that depends on that value, will multi execute twice instead of 4 times as in SME.

Moreover, MF further uses the view information provided by the encoding in order to multi execute less in case of branching commands. For example, for SME-TINI with SME memory µ(h) = (1 : 0 : 0 : 0) the program:

\[
\begin{align*}
1 & \quad \text{if } h = 0 \text{ then} \\
2 & \quad h := h + 1
\end{align*}
\]

executes 4 times (where the assignment at line 2 executes 3 times).

Using the MF memory encoding µ(h) = (⊤1 : 0 : 0 : 0), MF remembers that at line 2 there is no possible observation for the view ⊤ (because for view ⊤ the value of h is 1 so it doesn’t take the then branch). Hence, the assignment h := h + 1 only executes once with a memory where h is 0 (the view of variable h corresponding to the 3 levels which are not ⊤).

For a program h := l, where µ(l) is (1 : 1) in SME, MF keeps only the value 1: a single value represents the fact that all views can observe the same value. Thus the assignment h := l executes once (and all future executions dependent on h will also be reduced).
We present an extension to the original Multiple Facets mechanism [3] for an arbitrary security lattice \( (L, \sqsubseteq) \), which we call Generalised Multiple Facets mechanism, or GMF. Similarly to Multiple Facets, GMF operates over a faceted memory \( \mu \) that maps variables to simple values or faceted values. A faceted value is of the form \( \langle l ? V_1 : V_2 \rangle \) where \( l \in L \) is a security level, and \( V_1 \) can be either a faceted value or a simple value. The first facet \( V_1 \) of \( \langle l ? V_1 : V_2 \rangle \) is called private, and visible to the observers at security level \( l \) or higher levels in the lattice; the second facet \( V_2 \) is called public, and visible to security levels that are lower or incomparable to \( l \). We use \( V \) as a meta-variable for faceted values or simple values. Every evaluation in GMF (see Fig. 6) is marked with a set of security levels \( pc \), for which the current computation is visible.

### 3.1 Expression evaluation

By \( \hat{\mu}^{pc}(e) \) we denote the evaluation of expression \( e \) in faceted memory \( \mu \) with set of security levels \( pc \). The definition of \( \hat{\mu}^{pc}(e) \) is presented in Fig. 4. For example, consider the evaluation of \( x \) when the faceted value \( x \) in memory \( \mu \) is \( \langle l ? V_1 : V_2 \rangle \). To define which facet is useful given a \( pc \), we consider the following cases:

- All the levels in \( pc \) are greater than or equal to \( l \), denoted \( l \preceq pc \) (i.e. \( \forall l' \in pc. l \subseteq l' \)) - the evaluation can use the private facet \( V_1 \) because the public facet \( V_2 \) is anyway not useful for every level in this \( pc \).
- All the levels in \( pc \) are lower than or incomparable to \( l \), denoted \( l \not\preceq pc \) (i.e. \( \forall l' \in pc. l \not\subseteq l' \)) - the evaluation can only use the public facet \( V_2 \) because \( V_2 \) is a facet visible to any view that is lower than or incomparable to \( l \).
- Otherwise, we say that \( l \) and \( pc \) are incomparable and denote it by \( l \parallel pc \) (i.e. \( \exists l' \in pc. l \subseteq l' \land l \not\subseteq l' \)). We first evaluate \( V_1 \) with \( pc_1 = \{ l' \in pc \mid l \subseteq l' \} \) - the set of all levels in \( pc \) which are greater than or equal to \( l \). Then, we evaluate \( V_2 \) with \( pc_2 = pc \setminus pc_1 \) which is the set of all levels in \( pc \) which are lower than or incomparable to \( l \). Finally, we combine the two results in a new faceted value.

To evaluate a variable \( x \), we use a special unary operator \( \otimes^{pc}(\hat{\mu}(x)) \), which returns the value that is visible to all the levels in the \( pc \). Let’s consider the case of \( \otimes^{pc}(\langle l ? V_1 : V_2 \rangle) \). Notice that, if \( pc \) and \( l \) are incomparable, meaning that there are some levels in \( pc \) that are higher than or equal to \( l \) and other levels in \( pc \) that are lower than or incomparable to \( l \), denoted by \( l \parallel pc \), then the evaluation returns the faceted value \( \langle l \parallel pc_1 (V_1) \parallel \otimes^{pc_2} (V_2) \rangle \). The form of the result of \( \hat{\mu}^{pc}(e) \) is described in Lemma 3.1.

**Lemma 3.1.** If \( \hat{\mu}^{pc}(e) = \langle l ? V_1 : V_2 \rangle \), then \( l \parallel pc \).

**Example 3.2 (Expression evaluation).** Consider the lattice \( (L, \sqsubseteq) \) from Fig. 3, and the evaluation of \( x + y \) in \( \hat{\mu} \), where \( \hat{\mu}(x) = \langle M_1 : ? 10 : 0 \rangle \) and \( \hat{\mu}(y) = \langle M_2 : ? 5 : 0 \rangle \).

Suppose that \( pc = \{ M_1, H \} \). Since all the levels in \( pc \) are higher than or equal to \( M_1 \), the evaluation of \( x \) returns \( \hat{\mu}^{pc}(x) = 10 \). Since \( pc \) and \( M_2 \) are incomparable, the evaluation of \( y \) returns \( \hat{\mu}^{pc}(y) = \langle M_2 : ? 5 : 0 \rangle \). Next, the evaluation of \( 10 + \hat{\mu}^{pc}(y) = \langle M_2 : ? 5 : 0 \rangle \) is split into two: one uses a facet visible to \( M_2 \) (and hence \( H \)), and another one
\[ \begin{align*}
\hat{\mu}^{PC}(v) &= v \\
\hat{\mu}^{PC}(x) &= \emptyset^{PC}(\hat{\mu}(x)) \\
\hat{\mu}^{PC}(e_1 \oplus e_2) &= \hat{\mu}^{PC}(e_1) \oplus^{PC} \hat{\mu}^{PC}(e_2) \\
\emptyset^{PC}(v) &= v \\
\emptyset^{PC}(l \oplus v_1 : v_2) &= \begin{cases} \\
\emptyset^{PC}(v_1) & \text{if } l \leq pc \\
\emptyset^{PC}(v_2) & \text{if } l \not\leq pc \\
(l \oplus v_1 \oplus v_2) & \text{otherwise} \\
\end{cases} \\
v_1 \oplus^{PC} v_2 &= v_1 \oplus v_2
\end{align*} \]

\[ \begin{align*}
\hat{\mu}^{PC}(v_1 : v_2) &= \begin{cases} \\
v \oplus^{PC} v_1 & \text{if } l \leq pc \\
v \oplus^{PC} v_2 & \text{if } l \not\leq pc \\
(l \oplus v_1 \oplus v_2) & \text{otherwise} \\
\end{cases} \\
(l \oplus V_1 : V_2) &= \begin{cases} \\
V \oplus^{PC} V_1 & \text{if } l \leq pc \\
V \oplus^{PC} V_2 & \text{if } l \not\leq pc \\
(l \oplus V_1 \oplus V_2) & \text{otherwise} \\
\end{cases} \\
(l \oplus V_1 \oplus V_2) &= \begin{cases} \\
V \oplus^{PC} V_1 & \text{if } l \leq pc \\
V \oplus^{PC} V_2 & \text{if } l \not\leq pc \\
(l \oplus V_1 \oplus V_2) & \text{otherwise} \\
\end{cases}
\end{align*} \]

Figure 5: Functions for faceted and normal memories.

Figure 6: Multiple facets for arbitrary security lattice

\[ \begin{align*}
\hat{\mu}^{PC}(x + y) &= \hat{\mu}^{PC}(x) +^{PC} \hat{\mu}^{PC}(y) \\
&= \emptyset^{PC}((M_1 \oplus 10:0) +^{PC} \emptyset^{PC}((M_2 \oplus 5:0)) \\
&= 10 +^{PC} \langle M_2 \oplus 5:0 \rangle = \langle M_2 \oplus 15:10 \rangle
\end{align*} \]

3.2 Semantics

We abuse the notation and use \( I\) as a projection function on simple values, faceted values and faceted memories. For any \( V \), \( I(V) \) returns the value in \( V \) which is visible to users at level \( l \). For any \( \hat{\mu} \), \( I(\hat{\mu}) \) returns the memory in \( \hat{\mu} \) which is visible to users at level \( l \).

\[ I(v) = v \]

The projection function \( I \) is used in the definition of \( \hat{\mu}^{|\Gamma|} \) function that converts a faceted memory to a simple memory (see Fig. 5). The semantics of GMF is defined in Fig. 6 as a big-step evaluation relation \( \Gamma \vdash (P, \mu) \\ \emptyset^{GMF} \mu' \), where program \( P \) is executed in a memory \( \mu \) and a security environment \( \Gamma \) that maps variables to security levels in a given security lattice \( \langle L, \subseteq \rangle \).

The main rule GMF first constructs a faceted memory from the standard memory using the transformation \( \mu \uparrow^{def}_{\Gamma} \) from Fig. 5, where \( \emptyset^{GMF} \) is the greatest lower bound of \( L \). The resulting faceted memory keeps original value of each variable \( x \) in a private facet, and adds default values (defined by \( \text{def} \) function) in a public facet. In a special case when the level of \( x \) is the smallest level in a lattice, we keep only a simple value \( \mu(x) \) that is visible to all security levels. We then evaluate the program with the constructed faceted memory and \( pc = L \). The resulting faceted memory is transformed back to a normal memory by using the projection function \( \hat{\mu}^{|\Gamma|} \).

The semantics rules for skip, sequence and while loop are straightforward. The GAssign rule uses a faceted evaluation \( \hat{\mu}^{PC}(e) \) defined in Section 3.1.

Before describing the semantics of if instruction, we first define several auxiliary functions. Let \( \text{dom}(\hat{\mu}) \) be the domain of \( \hat{\mu} \) and \( y \) be a fresh variable, i.e. \( y \not\in \text{dom}(\hat{\mu}) \). By \( \hat{\mu} \uplus (y \mapsto V) \) we denote a new memory \( \hat{\mu}' \), such that \( \text{dom}(\hat{\mu}') = \text{dom}(\hat{\mu}) \uplus \{ y \mapsto V \} \) and for all \( x \in \text{dom}(\hat{\mu}) \), \( \hat{\mu}'(x) = \hat{\mu}(x) \). By \( \hat{\mu} \downarrow y \), we remove \( y \) from the domain of \( \hat{\mu} \), that is, \( \text{dom}(\hat{\mu}') = \text{dom}(\hat{\mu}) \downarrow \{ y \} \). The semantics rules for if instruction if \( e \text{ then } P_1 \text{ else } P_2 \) with \( \hat{\mu} \) and \( pc \). If \( e \) is evaluated to a constant value (tt or ff), then only \( P_1 \) or \( P_2 \) is evaluated (see rule GIf-C).

When \( e \) is evaluated to a faceted value \( l \oplus V_1 : V_2 \), we construct a new program if \( y \text{ then } P_1 \text{ else } P_2 \), where \( y \) is a fresh variable. From Lemma 3.1, we have that \( l \uplus pc \), and hence \( pc_1 = \{ l' \in pc \mid l' \subseteq l \} \) and \( pc_2 = pc \setminus pc_1 \) are non-empty. In this case, we run the new program if \( y \text{ then } P_1 \text{ else } P_2 \) twice: once with the “higher” view than \( l \), i.e., with \( pc_1 = \{ l' \in pc \mid l' \subseteq l \} \) and \( y \) set to a private facet \( V_1 \), and another time with “lower or incomparable view” than \( l \), i.e., with \( pc_2 = pc \setminus pc_1 \) and \( y \) set to a public facet \( V_2 \). We then combine the resulting memories using the \( \uplus \) operator. The combination of faceted memories is based on the fact that when \( pc \) is split into \( pc_1 \) and \( pc_2 \) in the GIf-S rule, all levels in \( pc_1 \) is larger than or equal to \( l \), and all levels in \( pc_2 \) is smaller than or incomparable to \( l \).
\[
\begin{array}{l}
\llbracket v \rrbracket = v \\
\llbracket (l \triangleright v_1 ; v_2) \rrbracket = \\
\quad \begin{cases}
\llbracket v_1 \rrbracket & \text{if } V_1 = V_2, \\
\llbracket (l \triangleright v_1 ; v_2) \rrbracket & \text{if } l_1 \subseteq l, l \subseteq l_2, V_1 = \llbracket l \triangleright v_1 ; v_2 \rrbracket, \\
\llbracket (l \triangleright v_1 ; v_2) \rrbracket & \text{if } l \subseteq l_2, V_2 = \llbracket l \triangleright v_1 ; v_2 \rrbracket, \\
\llbracket (l \triangleright v_1 ; v_2) \rrbracket & \text{otherwise}.
\end{cases}
\end{array}
\]

Figure 7: Optimisation of a faceted value.

Notice that the form of a faceted value constructed by combining values can be reduced. For example, a faceted value of the form \(\langle H ? M_{11} ; V_{12} \rangle \) can be reduced to \(\langle M_1 \triangleright V_{12} \rangle \) because \(M_1 \subseteq H\) and the projection of the original value at any level is either \(V_{11}\) or \(V_2\). We use the optimisation on the constructed faceted values from Fig. 7.

Therefore, in the GIF-S rule after the evaluation of \(P'\) in two contexts, we combine the resulting faceted memories \(\hat{\mu}'_1 \triangleright y\) and \(\hat{\mu}'_2 \triangleright y\) and apply an optimisation operator \(\llbracket \rrbracket\) for each newly constructed faceted value. The correctness of \(\llbracket \rrbracket\) used to optimize faceted values is proven in Lemma 3.3.

**Lemma 3.3.** For all \(l, \forall V, \text{ it follows that } l(V) = (l \llbracket V \rrbracket)\).

**Example 3.4 (Evaluation of if instruction).** Consider the security lattice \((L_0, \subseteq)\) from Fig. 3 and the evaluation of the following program \(P\) with \(pc = L_0\) and \(\hat{\mu}\), where \(\hat{\mu}(x) = \langle M_1 ? (H ? tt : ff) \rangle\).

1. \(\text{if } x \text{ then } z := 10 \text{ else } z := 5\).

The evaluation follows the GIF-S rule since \(M_1 \triangleright L_0\). We construct \(P' = \text{if } y_1 \text{ then } P_1 \text{ else } P_2\) and first evaluate \(P'\) with \(pc_1 = \{ l' \in pc | \{ M_1, H \} \text{ and } \hat{\mu}(l) = \mu(l) \Rightarrow (H ? tt : ff)\}\), and then evaluate \(P'\) with \(pc_2 = pc \setminus pc_1 = \{ M_2, L \}\) and \(\hat{\mu}(l) = \mu(l) \Rightarrow (y \Rightarrow tt)\).

Since \(pc_1 = \{ H, M_1 \}\) and \(\hat{\mu}^{pc}(y) = (H ? tt : ff)\), the evaluation of \(P'\) with \(pc_1\) and \(\hat{\mu}\) is split again to two evaluations: one with \(P'' = \text{if } y_1 \text{ then } P_1 \text{ else } P_2\) and \(pc_1 = \{ H \}\) and \(\hat{\mu}(l) = \hat{\mu}(l) \Rightarrow (H ? tt : tt)\); and the other one with \(P'' = \text{ if } y_1 \text{ then } P_1 \text{ else } P_2\) and \(pc_1 = \{ M_1, H\}\) and \(\hat{\mu}(l) = \hat{\mu}(l) \Rightarrow (H ? tt : ff)\).

The evaluation of \(P''\) with \(pc_1\) and \(\hat{\mu}\) follow the GIF-C rule and we get two faceted memories \(\hat{\mu}'_{11}\) and \(\hat{\mu}'_{12}\), where \(\hat{\mu}'_{11}(z) = 10\) and \(\hat{\mu}'_{12}(z) = 5\). Then, \(\hat{\mu}'_{11} \triangleright y_1\) and \(\hat{\mu}'_{12} \triangleright y_1\) are combined and we get \(\hat{\mu}'_{1}\), where \(\hat{\mu}'_{1}(z) = (H ? 10 : 5)\).

The evaluation of \(P_2\) with \(pc_2\) follows the GIF-C rule and the result is \(\hat{\mu}'_{2}\), where \(\hat{\mu}'_{2}(z) = 10\). At this point, \(\hat{\mu}'_{1} \triangleright y_1\) and \(\hat{\mu}'_{2} \triangleright y_1\) are combined and the result is \(\hat{\mu}'\), where \(\hat{\mu}'(z) = (M_1 ? (H ? 10 : 5) : 10)\).

**Example 3.5 (Evaluation with the GFM rule).** Consider the lattice \((L_0, \subseteq)\) from Fig. 3 and program \(P\) from Example 3.4 with one more instruction \(x := x_1 > x_2\). Suppose that \(\Gamma(x_1) = M_1\), \(\Gamma(x_2) = H\), \(\Gamma(z) = H\), \(\mu(x_1) = 10\), \(\mu(x_2) = 5\), the default values for \(x_1\) and \(x_2\) are respectively 100 and 20. Let \(\mu = \mu^{def}\). It follows that \(\hat{\rho}(x_1) = \langle M_1 ? (H ? tt : ff) \rangle\). We next establish the relation between the evaluation with GFM and the standard semantics.

Following GFM rule, the program is evaluated with \(pc = L_0\). For the assignment instruction, the value of \(x\) is updated to \(\hat{\rho}^{pc}(x_1 > x_2) = \langle M_1 ? (H ? tt : ff) \rangle\). The rest of the evaluation is described in Example 3.4, and the resultant faceted memory is \(\hat{\mu}'\), where \(\hat{\mu}'(z) = (M_1 ? (H ? 10 : 5) : 10)\).

The memory after the application of rule GFM is \(\hat{\mu}' = \hat{\mu}'\). Since \(\Gamma(z) = H\), the value of \(z\) is \(\hat{\mu}'(z) = H\langle M_1 ? (H ? 10 : 5) : 10 \rangle\).

**3.3 Equivalence to SME-TINI and Security Guarantee**

**SME-TINI.** The semantics of SME-TINI, termination-insensitive version of SME, for an arbitrary security lattice is presented below, where \(\hat{\mu}\) is a vector that maps levels to normal memories; \(\mu \triangleright l\) constructs a memory where \(\mu\) is a faceted memory at level \(l\) and \(\forall l, \Gamma\) constructs a memory by combining all memories in \(\hat{\mu}\); and \(\hat{\mu}'\) is a function mapping variables to default values.

**Example 3.6 (TINI).** An enforcement mechanism \(A\) is termination insensitive non-interferent (TINI) if for all security environments \(\Gamma\), programs \(P\), and memories \(\mu_1\) and \(\mu_2\), we have \(\forall l, \mu_1 \triangleright l \Rightarrow (P, \mu_1) \triangleright (P, \mu_2) \Rightarrow (P, \mu_2) \triangleright (P, \mu_1) \Rightarrow (P, \mu_2) \triangleright (P, \mu_1)\).

**Theorem 3.7.** SME-TINI is TINI.

Equivalence to SME-TINI. To prove the equivalence between GFM and SME-TINI, we formally define the semantic equivalence of two mechanisms.

**Definition 3.8.** Two enforcement mechanisms \(A\) and \(B\) are equivalent if for any \(\Gamma\) and \(\mu\), we have that \(\forall l, \mu \triangleright l \Rightarrow (P, \mu) \triangleright (P, \mu) \Rightarrow (P, \mu) \triangleright (P, \mu)\).

**Lemma 3.9.** \(\langle P, \mu \rangle \triangleright l_{pc} \hat{\mu}(l) = l(l(\hat{\mu}))\) for all \(l \in pc\).

Thanks to Lemma 3.9, we now prove the equivalence of GFM and SME-TINI.

**Theorem 3.10.** GMF and SME-TINI are equivalent.

As a consequence, we have that GMF is TINI.

**Remark 3.1.** MF [3] is constructed for a set of principals. When the set of principals is fixed, we can use GMF to encode MF: we construct the lattice \((P, \subseteq)\), where each element is a set of principals, we prove that GMF for \((P, \subseteq)\) and MF for \(P\) are equivalent [23].
4 OPTIMIZING GMF

In Section 3, we presented the semantics of Generalised Multiple Facets (GMF) for arbitrary lattice and have proven it to be equivalent to SME-TIN. However, GMF from Fig. 6 can be further optimised and avoid repeating evaluations of the same commands. The following example demonstrates the sub-optimality of GMF.

Example 4.1 (GMF is not optimal). We consider the below program from Example 2.1. The lattice is $\langle L_B, \sqsubseteq \rangle$ from Fig. 2.

1. \texttt{winner := 0;}
2. \texttt{test := (x1 ≤ x2) and (x2 ≤ x3);}
3. \texttt{if test then winner := 2 else skip}

Suppose that the bid offers of $B_1, B_2,$ and $B_3$ are respectively 10, 5, and 7, and the default values for $B_1$ are 0. W.r.t. this setting, the initial faceted memory is $\hat{\mu}$, where $\hat{\mu}(x_1) = (B_1 \uparrow 10 : 0), \hat{\mu}(x_2) = (B_2 \uparrow 5 : 0)$, and $\hat{\mu}(x_3) = (B_3 \uparrow 7 : 0)$. We consider the execution of the program with GMF.

After line 2, \texttt{test := (B1 ? Φ : Φ : (B2 ? Φ : tt : tt : tt))}. Following the semantics of GMF, the assignment winner := 2 is evaluated twice with $pc_B = \{B_3\}$, and $pc_{c_1} = \{\}$; the \texttt{skip} instruction is evaluated three times with $pc_T = \{\top\}$, $pc_B = \{B_1\}$, and $pc_B = \{B_2\}$.

The main idea of our optimisation lies in reducing the number of sub-evaluations and hence the number of faceted memory combinations. For Example 4.1, we propose a mechanism that merges the evaluations corresponding to $pc_B = \{B_3\}$ and $pc_c$ into one evaluation with $pc_c = \{B_3, \bot\}$. This simplification is possible since $test$ denotes the same value (i.e., $\top$) under $pc_B$ and $pc_c$. Similarly, our simplification merges the evaluations corresponding to $pc_T$, $pc_B$, and $pc_B$, where $test$ denotes $\bot$, into one evaluation with $pc_c = \{\top, B_1, B_2\}$, and thus evaluates each branch of the if command only once.

In this section, we propose semantics of optimised GMF (OGMF) that reduces the number of sub-evaluations, and hence is more resource-friendly than GMF.

4.1 Semantics

The ideas behind the OGMF rule, and the rules for skip, assignment, sequence, and while instructions and similar to the corresponding ones of GMF. The functions $\mu \uparrow^\text{def}_T (x)$ and $\hat{\mu}_1(x)$ are defined in Fig. 5. We now explain the semantic rules for the conditional instruction.

Consider evaluation of the program \texttt{if e then P1 else P2} with $pc$ and memory $\hat{\mu}$, and $\bar{\rho}^{pc}(e) = V$. In order to evaluate each branch of the conditional only once, we split the $pc$ in two subsets: in the first subset $pc_c$ the visible value of $V$ is true, and in the remaining subset $pc_c$, $V$ is false. We now have three distinct cases.

If $pc_c = pc$, meaning that for all levels in $pc$, the visible value of $V$ is true, then $P_1$ is evaluated (rule OIf-T). If $pc_c = pc$, then for all levels in $pc$, the visible value of $V$ is false, and only $P_2$ is evaluated (rule OIf-F). Finally, when $pc$ is split in non-empty $pc_c$ and $pc_c$, then both $P_1$ and $P_2$ are evaluated, and their results ($\hat{\mu}_1$ and $\hat{\mu}_2$) are combined by $\hat{\mu}_1 \oplus^{pc_c, pc_c} \hat{\mu}_2$ (rule OIf-S) to a new faceted memory. The intuition behind this combination is that the projection of $\hat{\mu}_1 \oplus^{pc_c, pc_c} \hat{\mu}_2$ at $l \in pc_c$ is taken from the evaluation of $P_1$ and its projection at $l \in pc_c$ is taken from the evaluation of $P_2$.

The definition of combination of memories for GMF (bottom of Fig. 8), we distinguish two cases. If for some variable $x$, its value in both faceted memories is the same, ($\hat{\mu}_1(x) = \hat{\mu}_2(x)$), then we do not need to construct a new faceted value. Instead, we optimise the current value using the optimisation operator from Fig. 7.

If the values of $x$ in $\hat{\mu}_1(x)$ and $\hat{\mu}_2(x)$ are different, then we construct a new faceted value $V = \mathbb{P}(V_1, V_2, pc_1, pc_2)$ and apply further optimisation on the resulting value $V$ using a new optimisation operator that takes into account a faceted value and the current $pc$: $\mathbb{P}[V, pc]$ optimizes the form of $V$ and is described in Fig. 9. We show an example of such optimisation in Example 4.4.

To combine two faceted memories, we first construct a new faceted value by using $\mathbb{P}(V_1, V_2, pc_1, pc_2)$:

$$\mathbb{P}(V_1, V_2, pc_1, pc_2) = \{\langle pc_1 \cup pc_2, V_1, V_2, pc_1, pc_2 \rangle\}$$

where $List(S)$ is a list of security levels from a set $S$, such that if $l$ appears before $l'$ in $List(S)$ then $l \not\sqsubseteq l'$. If the relation $\sqsubseteq$ in a given security lattice is not a total order, we can transform it into a total order $\sqsubseteq_T$ provided that $\sqsubseteq$ is a finite partial order. We can then view $List(S)$ as a list such that for any $l$ and $l'$ in this list, if $l$ appears before $l'$, then $l' \sqsubseteq_T l$.

The definition of $\mathbb{P}(V_1, V_2, pc_1, pc_2)$ uses the following operator that creates a faceted value based on an ordered list of security...
levels \( L \), two faceted values, \( pc_1 \) and \( pc_2 \):

\[
\langle L, V_1, V_2, pc_1, pc_2 \rangle = \begin{cases} 
\langle l(V_1) \rangle & \text{if } L = l, l \in pc_1, \\
\langle l(V_2) \rangle & \text{if } L = l, l \in pc_2, \\
\langle \langle l(V_1) \rangle \cap \langle l(V_2) \rangle, \langle pc_1 \cap pc_2 \rangle \rangle & \text{otherwise},
\end{cases}
\]

Notice that the form of the faceted value created by \( \langle V_1, V_2, pc_1, pc_2 \rangle \) may be suboptimal.

**Example 4.2 (Faceted value construction).** Suppose that \( V_1 = 2, V_2 = 0, pc_1 = \{ B_3, \perp \}, pc_2 = \{ T, B_1, B_2 \} \). If \( \langle l, v : V' \rangle \) has the form \( \langle l \rangle \) where \( V' \) is either a simple value or a faceted value.

If \( n \) is the form \( \langle l \rangle \), then the faceted value is straightforward. We now consider the case when \( V \) is of the form \( \langle l : \langle l' \rangle \rangle \) and \( v = v' \). For demonstration, consider the lattice \( \langle L, B, \subseteq \rangle \) from Fig. 2.

If the faceted value \( V \) is of the form \( \langle l : \langle l' \rangle \rangle \) then the optimisation is straightforward. We now consider the case when \( V \) is of the form \( \langle l : \langle l' \rangle \rangle \) and \( v = v' \). For demonstration, consider the lattice \( \langle L, B, \subseteq \rangle \) from Fig. 2.

If the faceted value \( V \) is of the form \( \langle l : \langle l' \rangle \rangle \) and \( v = v' \), and for all the levels in the \( pc \), for which either \( B_0 \) or \( B_2 \) is visible, it is guaranteed that they observe the same value \( v \) because \( cond(V, pc) \) below, then we distinguish the following two cases:

\[
cond(V, pc) \iff \forall l_i \in pc : glb(l, l_i) \subseteq l_i \implies l_i(V) = v.
\]

- If all levels in \( pc \) are greater than or equal to \( glb(l, l_i) \) (i.e., \( glb(l, l_i) \triangleq pc \)) then \( V \) is reduced to \( v \). For example, if \( pc = \{ B_1, B_2, B_3 \} \), \( glb(B_1, B_2) = \perp \), \( glb(B_1, B_3) = \perp \), and the \( cond(V, pc) \) we know that \( B_1(V) = B_2(V) = B_3(V) = v \) then we can reduce such faceted value to simply \( v \) because \( V \) is not useful for such \( pc \).

- If only some levels in \( pc \) are greater than or equal to \( glb(l, l_i) \) (i.e., \( glb(l, l_i) \triangleq pc \)), then \( V \) is reduced to \( \langle l(V) \rangle \) and this value is reduced further recursively. Consider that we add one more security level \( L \) to the lattice \( \langle L, B, \subseteq \rangle \) such that \( L \subseteq \perp \). If \( pc = \{ B_1, B_2, L \} \), \( glb(B_1, B_2) = \perp \), then \( glb(B_1, B_2) \triangleq pc \) because \( L \subseteq \perp \). We then construct a set of security levels \( S \) from \( pc \), which are higher or equal than \( glb(l, l_i) \), and therefore the view on \( V \) from these levels is \( v \) (because \( cond(V, pc) \) holds). In our example, \( S = \{ B_1, B_2 \} \), and we construct a new faceted value \( V' = \langle \langle L \rangle \rangle \) such that \( V' = \langle \langle L \rangle \rangle \triangleq L \). We then find a new \( pc' = (pc \triangleq S) \cup \{ glb(l, l_i) \} \triangleq (L, L) \), and we need to keep \( glb(l, l_i) \) in \( pc' \) because we must ensure that all the levels present in the new faceted value are also present in \( pc \). Therefore, the reduced faceted value for our example is \( \langle \langle l(V) \rangle \rangle \). Finally, if none of the above conditions hold then we recursively reduce the faceted \( V' \).

The correctness of \( \hat{\mu}_1 @ pc_1 \) in the OIf-S rule is proven in Lemma 4.3.

**Lemma 4.3.** For all levels \( l \), variables \( x \), sets of security levels \( pc_1 \) and \( pc_2 \), and memories \( \hat{\mu}_1 \) and \( \hat{\mu}_2 \).

- If \( l \notin pc_1 \), then \( l(\hat{\mu}_1 @ pc_1 \cap pc_2) = l(\hat{\mu}_1(x)) \).
- If \( l \in pc_2 \), then \( l(\hat{\mu}_1 @ pc_1 \cap pc_2) = l(\hat{\mu}_2(x)) \).

**Example 4.4 (Optimisation of faceted value).** Consider a faceted value \( \langle T \rangle : \langle B_1 \rangle : \langle B_2 \rangle : \langle B_3 \rangle \). We evaluate it with \( pc = \{ T, B_1, B_2, B_3, \perp \} \) from Example 4.2. We show how this value is optimised with our optimisation function \( \hat{\mu} \).

As a consequence, OGMF and GMF are equivalent even though OGMF is optimized. In addition, OGMF is TINI.

**4.2 Equivalence to SME-TINI and Security Guarantee**

We first establish the relation between the standard semantics and the semantics of OGMF.

**Lemma 4.6.** \( (P, \hat{\mu}) @ OGMF \) iff \( (P, l(\hat{\mu})) @ IGMF \) for all \( l \in pc \).

We now prove the semantic equivalence result for OGMF and SME-TINI.

**Theorem 4.7.** OGMF and SME-TINI are equivalent.

As a consequence, OGMF and GMF are equivalent even though OGMF is optimized. In addition, OGMF is TINI.
\[ ([l \downarrow v; v'], pc] = \begin{cases} v & \text{if } v = v' \\ ([l \downarrow v; v']) & \text{otherwise}. \end{cases} \]

\[ ([I \downarrow v; v', pc] = \begin{cases} ([I'] \downarrow v'; v'), pc' & \text{if } I' \subseteq I, v = v', \text{where } pc' = pc \setminus \{I\}, \\ v & \text{if } I \parallel I', v = v', \text{cond}(V, pc) \text{ and } \text{glb}(l, I') \prec pc, \\ ([\text{glb}(l, I') \downarrow v; v''), pc'] & \text{if } I \parallel I', v = v', \text{cond}(V, pc) \text{ and } \text{glb}(l, I') \parallel pc, \text{ where } pc' = (pc \setminus S) \cup \{\text{glb}(l, I')\}, \\ ([l \downarrow v; [l \downarrow v'; v'], pc] & \text{otherwise, where } pc' = pc \setminus \{I\}. \end{cases} \]

\[ \llangle (L, V) \rrangle = \begin{cases} (V) & \text{if } L = I, \\ ([l \downarrow V; \llangle T, V \rrangle) & \text{if } L = I, T, T \neq \{\}. \end{cases} \]

Figure 9: Definition of \([V, pc]\), and optimisation of a faceted value \(V\) with respect to the set of security levels \(pc\).

5 A TERMINATION SENSITIVE VERSION OF MULTIPLE FACETS

A termination sensitive model assumes that an attacker can observe termination of evaluations. In [19], the model is explained further: an attacker at level \(I\) can observe the termination of evaluations at level \(I\) and lower. In the case of GMF and OGMF, an evaluation marked with \(pc\) is an evaluation at \(I\) if \(I \in pc\). Notice that an evaluation is at more than one level whenever \(pc\) is not a singleton.

As illustrated by Example 5.1, GMF and OGMF do not prevent the influence of private data at higher levels to the termination of the evaluations at lower levels. In other words, GMF and OGMF do not prevent leakage on termination channel [19].

Example 5.1. Suppose that \(L = \{L, H\}\), where \(L \subseteq H\). We look at the evaluation of \(\text{if } x \text{ then (while } t \text{ do skip) else skip}\) with \(pc = L\) and \(\mu(x) = \{H? t : Hf\}\). When GMF or OGMF is used, the evaluation is split into two: one is with \(pc_1 = \{H\}\), the other one is with \(pc_2 = \{L\}\). The evaluation with \(pc_2\) converges, while the evaluation with \(pc_1\) diverges since its executing program is while \(t\) do skip. Therefore, the evaluation of the whole program with \(pc = \{L, H\}\) also diverges and hence, to an attacker at \(L\), the evaluation at \(L\) diverges. However, if the program is evaluated with \(\mu'(x) = \{H? t : Hf\}\), to the attacker at \(L\), the evaluation at \(L\) converges. Based on observations on those two evaluations, an attacker at \(L\) can gain insight about the high facet of \(x\). In other words, GMF and OGMF do not prevent the influence of data at \(H\) to the termination of the evaluation at \(L\).

Therefore, we propose Termination Sensitive Multiple Facets (TSMF), a version of MF that takes into account the termination sensitive model. TSMF is a generalization of a version of MF presented in [8, Appendix A]. The basic idea of TSMF is that when an if instruction is evaluated, TSMF performs a bounded evaluation of the instruction by using OGMF. If the OGMF evaluation does not terminate within the given time bound, then the instruction is evaluated instead using SME semantics with a low-prio scheduler [14]. The security guarantees offered by TSMF are the same as SME with the same low-prio scheduler [19]. The semantics of TSMF and the proofs about its security guarantees can be found in [23].

6 RELATED WORK

SME. Devriese and Piessens introduce the idea of Secure Multi-Execution [14]. Since then, many researchers have developed different aspects of this approach. Close to our work, Kashyap et al. [19] discuss how schedulers might affect security guarantees (i.e., TSNI and TINI) based on the chosen scheduler and the lattice ordering. They show several schedulers and classify them according to the strength of security guarantees and according to fairness properties.

This work complement theirs by providing a similar analysis but for an interplay of MF and SME semantics. SME [14] has many implementations: as a library in Haskell [18], as an experimental web browser based on Firefox [13], as a static program transformation for both Python and JavaScript [4], and as an adaptation to reactive systems [6]. In the work above, SME preserves the semantics of secure programs up to interleaving of events. To remedy that, Zararini et al. [37] carefully leverage SME to design a precise monitor which exactly preserves semantics of secure programs up to termination.

Several other works [10, 26, 33] expand SME and introduce declassification. In this work, we focus on semantics guarantees up to interleaving of events—as in the SME original formulation.

MF. Austin and Flanagan introduce MF semantics [3]—a technique often referred as an optimization for SME. However, as shown by Bielova and Rezk [7], they do not provide the same security guarantees (i.e., TINI vs. TSNI) and differ in their treatment of default values. This work provides yet another look into a comparison between both techniques to show their differences, while introducing novel value-based optimizations to MF. Another work by the same authors [9] compare and contrast five dynamic techniques, including MF and SME, to mainly reason about the preservation of semantics of secure programs, a property known as transparency. In this work, we show that GMF and OGMF enjoy the same transparency guarantees as SME-TINI (Theorems 3.10 and 4.7).

Tools. Most information flow control tools provide TINI, e.g., Jif [21], FlowCaml [25], Laminar [27], Paragon [11], and JSFlow [16]. Similarly, termination leaks are often ignored in security tools coming from the operating system research community, e.g., Asbestos [15], HiStar [38], and Flume [20]. A few exceptions to this trend are the security libraries LIO [31] and MAC [28], which provide TSNI for concurrent programs.

Decentralized label models. The decentralized label model (DLM), allows one to express the interests of mutually-distrusting principals without a central authority [22]. The set of labels forms a pre-order where the order relationship does not require to know all the points in the relationship to determine the result of comparing two labels—bearing in mind that there might be an infinite number of labels due to the dynamic creation of principals at runtime. In a similar spirit, DC-labels [32, 34] provides a decentralized label format which allows one to express rich policies dictated...
by mutually-distrusting principals as propositional logic formulas (without negation). In this work, we require to know all the points in the chosen lattice in order to optimize MF as shown by OGMF. Extending our techniques to DLM or DC-labels is an interesting direction for future work.

7 CONCLUSION AND PERSPECTIVES

This work contributes to develop techniques to secure programs using dynamic information flow—a promising approach to secure existing JavaScript code. We especially focus on proposing a technique that achieves a smaller number of executions than MF (and hence smaller memory footprint) without diminishing security guarantees. We further extend our MF-based technique to work with arbitrary finite lattices (GMF) based on the observation that

\[(or future ones) to work on potentially infinite lattices like the DLM\]

TSMF prioritizes resource usage as long as there are no risks for termination leaks. We expect that these insights will help inform future development of multi-execution-based techniques. In fact, an intriguing question is what it would take for our optimizations (or future ones) to work on potentially infinite lattices like the DLM or DC-labels—an interesting direction for future work.

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